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POLYNOMIAL MANIPULATION SYSTEM -FORTRAN IV PROGRAM

PREPARED BY

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#### FOREWORD

This report is a technical summary reporting the progress of a study conducted in the Mathematics Department and the Computer Center of Auburn University. The study is focused toward fulfillment of Contract No. DAAHO1-68-C-0296 granted to Auburn University by the Army Missile Command, Huntsville, Alabama.

#### ACKNOWLEDGEMENT

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#### ABSTRACT

A FORTRAN IV program which implements the Polynomial Manipulation System (PMS) is presented and described. PMS uses the Euclidean Algorithm to reduce a system of polynomials in several variables to a resultant system which can be solved sequentially as polynomials in one variable (Kronecker's method). PMS is described briefly and references are given to more complete discussions and to other pertinent literature.

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#### I. INTRODUCTION

The Polynomial Manipulation System (PMS) uses the Euclidean Algorithm for finding the eliminant and the greatest common divisor (g.c.d.) of two multi-variable polynomials. All polynomials involved are represented symbolically; PMS is a computer program whose input is the symbolic representation of two polynomials and whose output is (normally) the symbolic representation of their g.c.d. and eliminant.

The underlying theory and the application of PMS to the problem of solving systems of polynomial equations is discussed in [1], [2], and [3]. The present report describes a FORTRAN IV implementation of PMS developed on the IBM 360 Model 50 at Auburn University.

The program is described in Section II, the basic flow charts are given in Section III, Input/Output is discussed in Section IV, and efficiency of the method is discussed in Section V along with possible future work. The FORTRAN program is reproduced in Appendix A. Appendix B contains a simple example of the use of PMS for reduction of three polynomial equations in three variables to a resultant system which can be solved in sequence as polynomials in one variable.

#### II. PROGRAM DESCRIPTION

The PMS program is basically a main program with four subroutines, only one of which is significant. The other three subroutines are used for output, format headings on printed output and scaling of coefficients when they become large enough to possibly cause an overflow. In its present form the program is limited in that it is set up to use only 175K of IBM 360 storage. This limitation places constraints on the program which allows storage of only 50,000 polynomial entries (each term has n + 1 entries where n is the number of variables in the polynomial) which are presently set up as follows:

- 1) The pair of polynomials has, at most, four variables.
- 2) Each polynomial has at most 3160 terms.
- 3) The leading coefficient polynomials, to be defined below, can have, at most, 400 terms.

Minor modifications could increase the number of terms or variables or size of leading coefficients at the cost of decreasing the others or by use of a greater amount of machine storage. Still larger polynomials could be processed by use of tape, disc or other storage, but this has not been effected since such increases would only tend to accentuate certain disadvantages of the method to be discussed in Section V.

Consider the pair of polynomials U,T :  $E^N \rightarrow \mathbb{R}$ . Let  $\mathbf{x}_1, \dots, \mathbf{x}_n$  denote the variables. Each of these polynomial functions can be considered as a polynomial in  $\mathbf{x}_1$  whose coefficients

would then be polynomial functions from  $E^{N-1}$  to R. Let  $U^0$  and  $T^0$  denote the polynomials in  $x_2, \ldots, x_n$  which are the leading coefficients of U and T respectively considered as polynomials in  $x_1$ , and let u and t denote the degrees of U and T in  $x_1$ . We may assume  $t \ge u$ . Consider the polynomial R defined by

 $R = U^0T - T^0Ux_1^{t-u}$ 

R is a polynomial in  $x_1, \dots, x_n$ . Considering R as a polynomial in  $x_1$  with polynomial coefficients, it is seen that Degree(R)<t. Let Degree(R) = r. If  $r \ge u$ , let T = R and repeat above procedure. If r < u, let T = U and U = R and repeat the above procedure. After a finite number of applications of this algorithm a polynomial R will be found whose degree in  $x_1$  is zero. Thus R will be a polynomial in  $x_2, x_3, \dots, x_n$ . It is easily seen that at each stage R has any zeros that are common to U and T. The R which is free of  $x_1$  is called the eliminant of U and T.

# III. BASIC FLOW CHARTS

The flowcharts for output and scaling will be omitted as their detail is not significant to the main purpose of the program. The main program flow chart is given on page 4.

# MAIN PROGRAM

Read U,T

Compute U<sup>0</sup>, T<sup>0</sup>

Print U,T

 $\overrightarrow{1}$ 

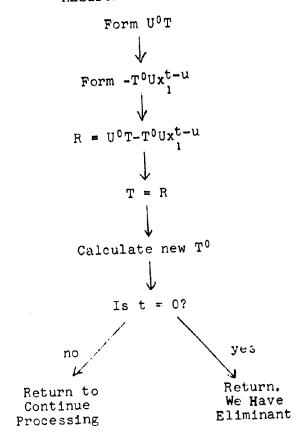
Determine which of U and T has greatest degree in x<sub>1</sub>. If it is T continue, if not interchange U and T, U and T.

Call RESIDUE and Form R

Scale if necessary

If R is free of  $x_1$ , print R and return to beginning to read two new polynomials. If not, let T = R.

# RESIDUE SUBROUTINE



#### IV. INPUT/OUTPUT

The polynomials U and T are read in separately. The first card for each polynomial contains the number of variables in the polynomial in columns 31-34 in integer format, right justified. The number of terms in the polynomial appear in columns 35-38 in integer format, right justified. Following this card the terms of the polynomial appear, one term per card. The coefficient appears in columns 1-16 in E format; following this are the exponents of the variables right justified in integer format in columns 17-21, 22-26, 27-31, etc.

The output of this program is available in any medium, although the program is currently set up for printed output only.

#### V. EFFICIENCY AND FUTURE WORK

The PMS program has not proved useful as a method of reducing polynomial systems of equations to a resultant system for the following reasons:

- 1) Storage efficiency is low. An inordinate amount of core is needed to process many simple appearing systems of equations.
- 2) Time efficiency is low. Extreme amounts of time are needed to solve all but the most simple problems. As few as four equations in four variables with small exponents (on the order of ten or less) take many hours of machine time to reach a solution. Simpler problems are solvable in

small amounts of time, but other methods without these disadvantages can be used to solve these systems.

3) Certain types of systems give solutions which have a low order of accuracy. Several articles, [4], [5], [6], have been published discussing this problem as well as the two above.

The basic PMS program will be examined and modified to determine if it is of value in algebraically solving simple systems of differential equations.

#### VI. REFERENCES

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  Variables for a Digital Computer, Duke University

  Memorandum 61-5-5,1 (1 May 1961).
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  Auburn University Report No. AU-T-2, The First Annual
  Report on Project THEMIS Research at Auburn University
  (October 31, 1968), 175-185.
- 4. Emmer, Edward J., Jr. <u>Multiple Solution Sets for Non-linear Mechanics Problems</u>, Ph.D. Thesis, Case Institute of Technology, 1966.
- 5. Collins, George T. <u>Subresultants and Reduced Polynomial</u>
  Remainder Sequences, J. ACM (Jan. 1967), 128-142.
- 6. Ku, S. Y. and Adler, R. J. <u>Computing Polynomial Resultants</u>: <u>Bezout's Determinant vs. Collins' Reduced P.R.S.</u>

  Algorithm, Comm. ACM (Jan. 1969), 23-30.

# NOT REPRODUCIBLE

#### APPENDIX A

# FORTRAM IV Program Listing

```
①农布农在交流布度条款布农布农布产业内部介有内容介绍产品介绍的介绍的产品产品产品产品产品介绍产品有产品有效的企业产品有效的产品有效的产品有效的企业。
       IMPLICIT INTO DEPRIEDO.

CORREDO HMADELES NOTOS CONTRACTORAS (2.400), TRANCACO), PERSON,
     NIR(4,2160), 179ATE, 1"1 PIC, ((1) a), (1) (A, P(0)), (( ) a), (T(A, P(0)),
     MI.S2, LS1, IM
       MAXP = 400
() 我们要在5个本有者或者或者或者或者或者的5、特别可能的约约中心中心中的4本有关。这个对于这种方式的专业的约约或者或者的成为的或者或者或者
【中的物域类的表现在在特殊的现在我们的中心,但是是一种我们的作用的自由各种的主要的专用实验和特别的实验和专业的基础的的特别的对抗性性
       330A40 =
       t_M±0.
       777 Mag = 3
       用设备((3+1) (16/4+ 27/27)
       FFC WARIALOW OF STORY
      151:0
      5.50:0
       ըց տալ,սու և դ
       PERF (5,47 G(1) + (10(1+1) + 1=1+NVAB1)
PERF (5,17 GV) - GF GV
       (M^{*}(t), t = 1, y \in [t, t] \times \mathbb{R}^{n}
       P(2)(5,4) = 1(4),(3'(1,4),1=1,00000)

P(2)(5,4) = 1,07,183
       R(J)=U(J)
       DEPT = L SWAF L
       13(1*4)=10(1*4)
       1FINVARI .Gr. "VARZI (4) TO 106
       MAX = MVAKZ
       MANAT = WAVEL + 1
       DO TOR A - CANAL GRAX
       10 102 C = 1, HIPOL
       544 Jam 1 10
       MV ssl = MAX
       MAX = NVARI
       Dewell = 0
       DO SO J =1, STORMI
       IF(JPROU - IR(1,3)) 49,50,50
49
       J9W9H = JR(1,J)
CO
       CHALLMAN
       DO 55 J = 1, NTERMI
       TF(JPSSH - IS(1,J)) 55,54,55
54
       JU = JU + 1
       (UE) \times (UE) \times MU
       IF (ABS (UMAX(JB)).CL.1.86) LM=1
       DO 53 " = 2.MAX
```

8

```
53
       IUMAX(K-1,JU) = IR(K,J)
55 --- CONTINUE
       00 10 J = 1.NTERM1
      - Ut J) = R(J)
       DOIOI=1, MAX
10(1,J)=IR(I,J)
           9 J = 1.NTERM2.
       00
    --- RfJ)=T(J)
       D091=1, NVAR2
       [[]]][=([,]]]]
       CALL PRINT(2,NTERM2,0,0,NVAR2)
     --- IFINVAR2 .GE. NVAR1) GO TO 1075
       MAX = NVAR1
       NVAR2 = NVAR2 + 1
       DO 104 J = NVAR2 ,MAX
       DO 104 K = 1.NTERM2
104
       IR(J,K) = 0
       NVAR2 = MAX
       MAX = NVAR2
1075
       JPWRT = ()
       DO 60 J =1.NTERM2
       IF(JPWRT - IR(1,J1) 51,60,50
       JPWRT = IR(1,J)
51
       CONTINUE
60
       0 = TL
       DO 56 J = 1, NTERM2
       IF(JPWRT - IR(1,J)) 56,57,56
JT = JT + 1
57
       TMAX(JT) = R(J)
       IF(ABS(TMAX(JT)).GE.1.E6) LM=1
       DO 58 K = 2.MAX
58
       ITMAX(K-1,JT) = IR(K,J)
56
       CONTINUE
       00 R J = 1.NTERM2
       T(J)=R(J)
       DD8I=1,MAX
R
       IT(1,J)=IR(1,J)
107
       IF(JPWRT-JPWRU) 70,71,71 -
70
       NN = JPWRT
       JPWRT = JPWRU
       JPWRU = NN
       MAXT=NTERM1
       IF(NTERM2.GT.NTERM1) MAXT=NTERM2
       D080[=1,MAX
       TEMP=U(I)
       U([]=T(])
       T(I)=TEMP
       DOBOU=1, MAXT
       TEMP=IU(J, I)
       1U(J,I)=IT(J,I)
80
       IT(J,1)=TEMP
       NN = JU
       JU = JT
       JT = NN
       NN = NTERM1
       NTERM1 = NTERM2
       NTERM2 = NN
       NN = NVAR1
```

```
NVAR1 = NVAR2
       NVAR2 = NN
       JJ=JU
      TF(JT.GT.JJ) JJ=JT
       MA = MAX-1
     00.73 J = 1, JJ
       (L)XAMU=WT
       (L)XAMT = (L)XAMU
       WT = \{L\}XAMT
       D073K=1,MA
       ITT=IUMAX(K,J)
      \neg \text{IUMAX}(K,J) = \text{ITMAX}(K,J)
73
       ITMAX(K,J)=ITT
71
     ··· CONTINUE
Ç
     * IF(LM.EQ.1) CALL SCALE2(NTERM1,NTERM2,MAX)
       C=14
  101 CALL PESTOU (NTERM1, NTERM2, MAX, JOWRU, JPWRT, JU, JT)
       IF(LM .FQ. 1) CALL SCALE2(NTERM1, NTERM2, MAX)
       IF(LM.FQ.1) GOTO 3001
       IF(LS2.EQ.1) CALL SCALE2(NTERM1, NTHKM2, MAX)
       £52=0
3001
       LM=0
300
       IF(LS1.E0.0) GBT0 107
       L$1=0
       (M=C
  301 CALL PRINT ( 3, JU, NTERMI, NTERM2, MAX)
       WRITE(6,10000) ITIMES
10000
      FORMAT(1H , 5HSCALE, 13)
      GD TO 100
    1 FORMAT(30X,214)
       FORMAT(E16.7,1015)
      END
```

```
C
       SUBROUTINE FOR PRINTING THE TWO POLYNOMIALS AND THE RESIDUE
C
C
     SUBROUTINE PRINT (L. JU, NTERM1, NTERM2, N)
IMPLICIT INTEGER * 2(I-N)
     COMM-IN I 'AX(400), IUMAX(3,400), ITMAX(3,400), TMAX(400), R(3160),
   NIR(4,3160), TTRATE, ITTMES, U(3160), IU(4,3160), T(3160), TT(4,3160),
   NES2, ES1, E4
IF(L.E0.3) #MITE(6,310)
     IF(L.FQ.2) APITE(6,311)
     IF(L.FO.1) WRITE(6,3)2)
  50 CALL PRINTH(II)
     00315K=1.Ju
     PRITH[6,93 ) R(K),(IR(I,K),I=1,V)
315
     CONTINH
    K = M-1
503
     COMPTMUS
     RETURN
  43 FORMAT(1X,F16.7,3X,10(15,5X))
311
     FORMATCIHI, THE POLYNOMIAL U. IST)
     FORMAT(14-, THE POLYNOMIAL TO IS!)
312
     FORMAT() :- , THE ELIMINANT (S.)
310
 540 FORMAT (EDX, 214)
 541 F 1RMAT(E16.7,1015)
    END
```

```
SUBSCUTINE FOR PRINTING HEADINGS
      SUBRUUTINE PRINTH(K)
C
                   THIS SUBROUTINE MERELY PRINTS COLUMN HEADINGS FOR THE
C
                   VIRIABLES DEPENDING ON THE NUMBER OF VARIABLES.
                   THAT IS ITS ONLY PURPOSE.
                                               IT WILL HANDLE UP TO 10
                   VARIABLES.
ſ,
      GO TO (21,22,23,24,25,26,27,28,29,30),K
21
       WR ITE(6,31)
      RETURN
13
       WRITE(6:32)
      RETHIRN
23
       WRITE(6,331
      RETURN
       WR ITE(6,34)
24
      RETURN
       WRITE(6,35)
      RETURN
       WRITE(6, 36)
      RETURN
       WR. ITE(6,37)
      RETURN
29
       1211E(6.-9)
      RETURN
23
       WR | TE(6, 39)
      RETURN
30
       WR IT(36,40)
      RET! TN
   31 FORMAT(1HO, )1HCOEFFICIENT, 10X, 4HX(1))
   32 FORMAT(1HO,11HCDEFFICIENT,10X,4HX(1),5%,4HX(2))
   33 FORMAT(1H0,11HCDEFFICIENT,10X,4HX(1),5X,4HX(2),5X,4HX(3))
   34 FORMAT(1H0,11HCOEFFICIENT,10X,4HX(1),5X,4HX(2),5X,4HX(3),5X,4HX(4)
     1)
   35 FORMAT(1H0,11HCOEFF[C1FNT,10X,4HX(1),5X,4HX(2),5X,4HX(3),5X,4HX(4)
     1.5X,4HX(5))
   36 FORMAT(1H0,11HCOFFFICIENT,10X,4HX(1),5X,4HX(2),5X,4HX(3),5X,4HX(4)
     1,5X,4HX(5),5X,4HX(6)}
   37 FDRMAT(1H0,11HCOEFFICIENT,10X,4HX(1),5X,4HX(2),5X,4HX(3),5X,4HX(4)
     1,5X,4HX(5),5X,4HX(6),5X,4HX(7))
   38 FORMAT(1H0,11HCOEFFICIENT,10X,4HX(1),5X,4HX(2),5X,4HX(3),5X,4HX(4)
     1,5X,4HX(5),5X,4HX(6),5X,4HX(7),5X,4HX(8))
   39 FORMAT(1H0,11HCOEFFICIENT,10X,4HX(1),5X,4HX(2),5X,4HX(3),5X,4HX(4)
     1,5X,4HX(5),5X,4HX(6),5X,4HX(7),5X,4HX(8),5X,4HX(9))
   40 FORMAT(1H0,11HCOEFFICIENT,10X,4HX(1),5X,4HX(2),5X,4HX(3),5X,4HX(4)
```

1,5X,4HX(5),5X,4HX(6),5X,4HX(7),5X,4HX(8),5X,4HX(9),5X,5HX(10))

END

```
FORMS RISTINE AND SPETS TERMS
        RESTOUT SUBROUTING
C
     SUBROUTING RESIDUINTU, NTT, NV, JPU, JPT, JU, JT)
C
      IMPLICIT INTEGER#2(I-N)
      COMMUNE UMAX(400), IUMAX(3,400), ITMAX(3,400), TMAX(400), R(3160),
    NIR(4, 3160), ITRATE, ITIMES, U(3160), IU(4, 3160), T(3160), IT(4, 3160),
    NUSZ, USI, UM
      DIMENSION
                13(5)
      MAXR=3160
      CCA = 9XAM
      THITGER SWITCH
C
      T = 1
      99 20 J = 1,8TT
      A=T(J)
      DO5210 K=1,NV
      IR(K)=IT(K,J)
5210
       IE(IE(1).GE.UPT)90T020
      [F(A38(A) - 1.56) 45,46,46
      L S2=1
46
      RETURN
      00 19 E = 1, JU
      R(I) = UVAX(I) *A
       \mathsf{TR}(1,1) = \mathsf{JR}(1)
      90.10 \text{ KK} = 2.00
       IR(KY,T) = IUMAX(KK-1,L) + IB(KK)
10
       44=1-1
       IF(MM.EQ.O)GUTD19
      DOBOKK=1,MM
       IF(IR(1,KK).NE.IR(1,I))GOTO80
      DOTOOKKK=2,NV
       IF(IR(KKK,KK).NE.IR(KKK,I))60T080
700
      CONTINUE
       R(KK)=R(KK)+R(I)
       1F(ABS(R(KK)).GT.1.E-9)GDT0702
       NINM=MM-1
       DOSODILK=KK.NNM
       R(ILK)=R(ILK+1)
       DU800NEO=1.NA
800
       IR(MPQ, IIK)=IR(NPC, ILK+1)
       I = I - I
702
       I = f - 1
       GOTIO 701
       CONTINUE
80
701
       IF(I.LT.DAXR)SUTU19
11
       WRIT=(6,12)
       FORMAT(191,25HT00 MANY TERMS IN RESIDUE)
12
       STIP
19
       I = I + I
       CONTINUE
20
       an 40 J = 1,710
       1=11(3)
```

13

```
nn5211 K=1+N√
       IR(*)=[11(K,J)
5211
       IF(I8(1).35.JPU)GOT040
       1F(ABS(A) - 1.F6) 47,48,48
       LS2=1
49
       RETURN
        nn 41 L = 1. JT
47
       A*(J)XANT- = (I)XA
       IR(1,1) = IR(1) + JPT - JPU
       pn: 30 K = 2+NV
        1R(K,T) = ITMAX(K-1,L) + IB(K)
30
        M4=1-1
        99380KK=1+MM
        1F(19(1,KK).Nb.IR(1,I))GOT0390
        DD3700KKK=2.NV
        IF(IR(KKK, KK).NE.IR(KKK,I))G010389
        CONTINUE
3700
        D (KK) = R (KK) + R(I)
        15 (ABS(R(KK)).GT.1.E-9) GOTO3702
        MMM=MM-1
        njakgall K=KK,NNM
        @([[K+1] =e([[K+1]
                                        NOT REPRODUCIBLE
        753070490=1+4V
        TM (NOB, IUK)=IR(RPQ,IUK+1)
 3800
        ! = ! - 1
        t = t - 1
 3702
        60103701
 240
        CHETIMUS
        16(I.LT. MAXE )SETUAL
 3771
         A (7011)
1 = 1 + 1
 41
         CONTINUE
 40
         1 - 1 - 1
         Jesse T=0
 99
         ( in == ( )
         90 105 M = 1.T
 100
         IR(IR(1,%) - JPERT) 105,105,102
         JPHET = IP(1,4)
 102
         CONTINUE
 105
         JPT = JPWAT
         10 ( (PURT) 130,142,130
 16.
         1.51 11
         301 = I
         SETERM
         0 = TU
         99711=1,400
         .0=(11)xAMT
         90733=1,3
         1π55 X(JJ, 11) =0
         0056J=1,1
  25
          1: (JPWRT - TR(1,J)) 56,57,56
         J = JT + 1
          (L)^{\circ} = (TL)yAvT
          18 (TMAX(JT).GF.1.F6) LM=1
          mmed K=2,NV
          TTMAX(K-1,J^T) = IP(K,J)
  r, :
          CONTINUE
          CONTINUE.
                                  14
```

```
SUPROUTINE SCALEZ(NTERMI, NTERMZ, NV)
        IMPLICIT 'NTEGER +2(I-N)
       COMMON UMAX(400), IUMAX(3,400), ITMAX(3,400), TMAX(400), R(3160),
     NIR(4,3160), ITRATE, ITIMES, U(3160), IU(4,3160), T(3160), IT(4,3160),
     NLS2+LS1+LM
       EQUIVALENCE (NVARI, MAX)
       MVX = M\Lambda
       NVAR 2=NVAR 1
       90 \cdot 1 \cdot 1 = 1.490
       UMAX(I) = UMAX(I)/1000.
       TMAX(I) = TMAX(I)/1000.
1
       DD = 7 \cdot J = 1, NTERM1
7
       U(J) = U(J)/1000.
       00 9 J = 1,NTERM2
       T(J) = T(J)/1000.
9
       ITIMES = ITIMES +
       RETURN
       FND
```

# APPENDIX B

This appendix presents an example problem. The following system of three polynomials in three variables is reduced to a single polynomial in one variable:

- (1)  $x_1 + x_2 + x_3 = 0$
- (2)  $x_1 + x_2^2 = 0$
- (3)  $x_1 + x_2^2 = 0$

First, x is eliminated between (1) and (2) producing the following printout which has been labeled for expository convenience:

COEFFICIENT 0.1000000E 0.1000000E 0.1000000E	01 01 01	*(1) 0 0	x(2) 0 1 0	<b>x(3)</b> 0 0 1	(1)
COEFFICIENT 0.1000000E 0.1000000E	01 01	*(1) 1 0	<b>x(</b> 2) 0 2	<b>x</b> (3) 0	(2)
COEFFICIENT 0.1000000E -0.1000000E	01 01 01	x(1) 0 0	x(2) 2 1	<b>x(3)</b> · 0	( <b>E</b> <sub>1</sub> )

Second,  $x_1$  is eliminated between (3) and (2) producing the following printout:

Finally (E<sub>1</sub>) and (E<sub>2</sub>) are treated as a pair of polynomials in two variables  $x_1$  and  $x_2$ . Then  $x_1$  ( $x_2$  in our first system) is eliminated, producing the following printout:

COEFFICIENT 0.1000000E -0.1000000E -0.1000000E	01 01 01	x(1) 1 0	x(2) 0 1	(E <sub>1</sub> )
COEFFICIENT -0.1000000E 0.1000000E	01 01	x(1) 0	x(5)	(E <sub>2</sub> )
0.100000E -0.200000E	01 01	x(1) 0	x(2)	(E <sub>3</sub> )

 $(E_3)$  is our eliminant free of x and  $x_2$ , so it can be solved. Using its solution  $(E_2)$  can then be solved. Using this (2) can be solved and the solutions to the resultant system (2),  $(E_2)$ ,  $(E_3)$  are the solutions to the system (1), (2), (3). Three other equations from these six could have been taken to form the resultant system.

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# APPENDIX C

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A FORTRAN IV program which implements the Polynomial Manipulation System (PMS) is presented and described. PMS uses the Euclidean Algorithm to reduce a system of polynomials in several variables to a resultant system which can be solved sequentially as polynomials in one variable (Kronecker's method). PMS is described briefly and references are given to more complete discussions and to other pertinent literature.					

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